Note on Transportation Problem with New Method for Resolution of Degeneracy

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Abstract  Transportation algorithm minimizes the cost of transporting goods from m origins to n destinations along m*n direct routes from origin to destination. There are three methods to solve transportation problem and their solution can be further improved using Modified Distribution Method (MODI Method). This method not only checks that the solution obtained by three methods is optimal but also modify the distribution to make it optimal. The only hiccup with this method is the case of degeneracy i.e. when number of allocations less than m+n-1. In this case, we have to make the number of allocations equal to m+n-1. To remove degeneracy we allocate some cells with dummy allocations, whose value is nearly equal to zero. But, till now there is no specific rule for allocating these cells. In this study, specific method has been proposed to remove this problem of degeneracy. Thus this study helps to remove major bottleneck from transportation algorithm.

Keywords  Transportation Problem, Degeneracy, Modified Distribution Method, Vogel Method, Least Cost Method

1. Introduction

Transportation theory is a name given to the study of optimal transportation and allocation of resources. The problem was formalized by the French mathematician Gaspard Monge(1871). In the 1920s A.N. Tolstoi was one of the first to study the transportation problem mathematically. In 1930, in the collection Transportation Planning Volume I for the National Commissariat of Transportation of the Soviet Union, he published a paper Methods of Finding the Minimal Kilometrage in Cargo-transportation in space. Tolstoi(1939) illuminated his approach by applications to the transportation of salt, cement, and other cargo between sources and destinations along the railway network of the Soviet Union. In particular, a, for that large-scale, instance of the transportation problem was solved to optimality.

Major advances in transportation theory were made in the field during World War II by the Soviet/Russian mathematician and economist Leonid Kantorovich Consequently, the problem as it is stated is sometimes known as the Monge–Kantorovich transportation problem. Kantorovich won the Nobel prize for economics in 1975 for his work on the optimal allocation of scarce resources, the only winner of the prestigious award to come from the USSR.

F.L. Hitchcock (1941) worked on the distribution of a production from several sources to numerous localities. Koopman also worked on the optimum utilization of transportation system and used model of transportation, in activity analysis of production and allocation.


M. Klein (1967) developed a primal method for minimal cash flows with applications to the assignment and transportation problems. Hadley(1972) also included transportation problem in his book: Linear Programming.

Lee (1972) and Ignizio(1976) used goal programming to solve transportation problem.


Kwak & Schniederjans(1985) framed goal programming solutions to transportation problem with variable supply and demand requirement. R.K. Ahuja (1986) developed an algorithm for minimax transportation problem. In the same Romero has done a survey of generalized goal programming also Currin worked on the transportation problem with inadmissible routes.
Romero (1991) has written a book on critical issues in goal programming, followed by Tamiz & Jones (1995) who has done a review of goal programming and its applications. Hemaida & Kwak (1994) developed a linear goal programming model for transshipment problem with flexible supply and demand constraints. Sharma et al. (1999) analyzed various applications of multi-objective programming in MS/OR.


In this study, for resolution of degeneracy method has been proposed. The problem of degeneracy can occur in initial solution or it may arise in some subsequent iterations.

Transportation Problem

- The transportation Problem is one of the special type of Linear Programming Problem in which objective is to transport various quantities of single identical goods that are initially stored at various origins, to different destinations in such a way that the total transportation cost is least.
- Vogel’s Approximation Method is one of the most efficient methods to obtain basic feasible solution, which is most near to the optimal solution.
- Then usually we check the optimality of the solution with the help of either stepping stone method (Charnes and Cooper, 1954) or Modified Distribution (MODI) method.

In order to check the optimality of the cost, we apply the optimality criteria: There must be \( (m+n-1) \) number of allocated cells (m-number of rows, n-number of columns). If the number of allocated cells is less than \( (m+n-1) \), then it is a degenerate Transportation Problem.

In this article we will propose a technique of resolving degeneracy, which is not given in textbooks.

There are three methods of obtaining a feasible solution:

1. North West Corner Method (Layman’s Method)
2. Least Cost Method (Business Man’s Method)
3. Vogel’s Approximation Method (Operations Method)

Formulation of a General Transportation Problem

Consider that there are \( m \) sources Source1, Source2, ………, Source \( m \) with capacities \( a_1, a_2, \ldots, a_m \) and \( n \)-destinations (sinks) with requirements \( b_1, b_2, \ldots, b_n \) respectively. The transportation cost from \( i \)th source to \( j \)th destination is \( c_{ij} \) and the amount shipped is \( x_{ij} \). If the total capacity of all sources is equal to the total requirement of all the destinations, what must be the values of \( x_{ij} \) with \( i=1,2,\ldots,n \) for the total transportation cost to be minimum?

<table>
<thead>
<tr>
<th>Source/ Destination</th>
<th>D1</th>
<th>D2</th>
<th>..</th>
<th>Dj</th>
<th>..</th>
<th>Dn</th>
<th>Availability</th>
</tr>
</thead>
<tbody>
<tr>
<td>S1</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>( a_1 )</td>
</tr>
<tr>
<td>S2</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>( a_2 )</td>
</tr>
<tr>
<td>\vdots</td>
<td></td>
<td></td>
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<td></td>
<td></td>
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<tr>
<td>Si</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>( a_j )</td>
</tr>
<tr>
<td>\vdots</td>
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<td></td>
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<td></td>
<td></td>
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</tr>
<tr>
<td>Sm</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>( a_m )</td>
</tr>
<tr>
<td>Requirements</td>
<td>( b_1 )</td>
<td>( b_2 )</td>
<td>..</td>
<td>( b_j )</td>
<td>..</td>
<td>( b_n )</td>
<td>( \sum a_i )</td>
</tr>
</tbody>
</table>
The Mathematical formulation of the transportation problem is given by

Minimize \[ Z = \sum_{i=1}^{m} \sum_{j=1}^{n} c_{ij} x_{ij} \]

Subject to

\[ \sum_{j=1}^{n} x_{ij} = a_i, i = 1, 2, \ldots, m \]

\[ \sum_{i=1}^{m} x_{ij} = b_j, i = 1, 2, \ldots, n \]

and \[ x_{ij} \geq 0 \] for all \( i \) and \( j \).

There are three methods to obtain the feasible solution

a) North West Corner Method
b) Least Cost Method
c) Vogel’s Approximation Method

\textbf{a) North West Corner Method}

This method is also called layman’s method, as in this method the only motive is to have a balance between demand and supply. This is the oldest method known and had been in use, even before any of the method is known. We have to follow the following steps to solve the problem using this method.

\textit{Algorithm}

Step 1: Locate the cell in the north-west (upper left) corner of the matrix of the data completely ignoring the transportation cost.

Step 2: Transport the minimum of demand and supply values with respect to that cell and subtract this minimum from the supply and demand values.

Step 3: Check whether exactly one of the row/column corresponding to the north-west corner cell has now zero supply/demand respectively or both of them has zero.

Step 4: Delete the row or column or both depending upon whichever is having zero

Step 5: Repeat the steps 1-4 with the reduced matrix unless all demand and supply becomes zero

The gist of this method is to start transporting from first source to first factory and then whichever is possible i.e., either from first source to second destination or from source 2 to first destination depending upon the demand and supply. The process continues unless all the units are transported.

\textbf{Problem 1}

The Oberoi car company has warehouses in Amritsar (\( W_1 \)), Nanded (\( W_2 \)) and Patna (\( W_3 \)) and markets in Noida (\( M_1 \)), Bangalore (\( M_2 \)), Kolkata (\( M_3 \)) and Gandhinagar (\( M_4 \))

\begin{tabular}{|c|c|c|c|c|c|}
\hline
\textbf{Factory} & \textbf{Markets} & \textbf{M_1} & \textbf{M_2} & \textbf{M_3} & \textbf{M_4} \\
\hline
W_1 & 24 & 23 & 43 & 14 & 60 \\
W_2 & 34 & 32 & 12 & 40 & 80 \\
W_3 & 20 & 22 & 23 & 41 & 100 \\
\hline
\textbf{Demand} & 40 & 70 & 80 & 50 & Total=240 \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|c|c|}
\hline
\textbf{Factory} & \textbf{Markets} & \textbf{M_1} & \textbf{M_2} & \textbf{M_3} & \textbf{M_4} \\
\hline
W_1 & 24(40) & 23 & 43 & 14 & 20 \\
W_2 & 32 & 12 & 40 & 80 \\
W_3 & 22 & 23 & 41 & 100 \\
\hline
\textbf{Demand} & 70 & 80 & 50 & Total=240 \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|c|c|}
\hline
\textbf{Factory} & \textbf{Markets} & \textbf{M_1} & \textbf{M_2} & \textbf{M_3} & \textbf{M_4} \\
\hline
W_1 & 24(40) & 23(20) & & & \\
W_2 & 32 & 12 & 40 & 80 \\
W_3 & 22 & 23 & 41 & 100 \\
\hline
\textbf{Demand} & 50 & 80 & 50 & Total=240 \\
\hline
\end{tabular}

The cost associated with this basic feasible solution is computed as follows

\[ 24*40+23*20+32*50+12*30+23*50+41*50=6580 \]

\textbf{b) Least cost Method}

This method is also called Business man’s method, as in this case a business that is having no knowledge of operations management is considered. His only motive to reduce the transportation cost, while having a balance between demand and supply. He always tries to allocate items to the cells having minimum transportation cost.

\textit{Algorithm}

Step 1: find the cell with minimum per unit transportation cost.

Step 2: Transport the minimum of demand and supply values with respect to that cell and subtract this minimum from the supply and demand values.

Step 3: Check whether exactly one of the row/column corresponding to the north-west corner cell has now zero
supply/demand respectively or both of them has zero.
Step 4: Delete the row or column or both depending upon whichever is having zero
Step5: Repeat the steps 1-4 with the reduced matrix unless all demand and supply becomes zero
Consider Oberoi Car Company problem again using least cost method.
Tables below explain the Least Cost Method with problem1

Table 1: Least Cost Method

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M_1</th>
<th>M_2</th>
<th>M_3</th>
<th>M_4</th>
<th>Availability</th>
</tr>
</thead>
<tbody>
<tr>
<td>W_1</td>
<td>24</td>
<td>23</td>
<td>43</td>
<td>14</td>
<td>60</td>
</tr>
<tr>
<td>W_2</td>
<td>34</td>
<td>32</td>
<td>12(80)</td>
<td>40</td>
<td>80</td>
</tr>
<tr>
<td>W_3</td>
<td>20</td>
<td>22</td>
<td>23</td>
<td>41</td>
<td>100</td>
</tr>
</tbody>
</table>

Demand 40 70 80 40 Total=240

The total transportation cost is
23*10+14*50+12*80+20*40+22*60=4010
As it can be observed that the total cost obtained by this method is very less than the total cost obtained by North West Corner method. The major disadvantage with North West Corner method is that cost does not play any role in that method. The only aim of using that method is to a balance between demand and supply.

The disadvantage with this method is that finally we have to transport some of the items with very high cost.

c) Vogel’s Approximation Method (VAM)

This method is also known as penalty method. This method is based upon the concept that if you miss best (lowest cost) in first attempt you have to face the penalty in next attempt. Hence, for each source and destination, we compute a ‘penalty’ rating which is the difference in cost of the two cheapest routes for the source and destinations. The major advantage of this method is solution given by this method is closer to the optimal solution.

Algorithm

Step 1: From the transportation table, we determine the penalty for each row and column. The penalties are calculated for each row (or column) by subtracting the lowest cost element in that row (column) from the next lowest cost element in the same row (column). Write down the penalties below the rows (aside the columns) of the table.

Step 2: Select the row (column) with the highest penalty rating and allocate as much as possible from the supply and requirement values to the cell having the minimum cost. If there is a tie in the values of penalties, then check the next level penalty i.e., the next two minimum cost, the route having more next penalty will be chosen. (New Rule for a tie of penalties)

This rule is extremely helpful in obtaining basic feasible solution in close proximity to the optimal solution.

Step 3: Adjust the supply and demand conditions for that cell. Eliminate those rows (columns) for which the supply and demand requirements are met.

Step 4: Repeat the steps 1-3 with the reduced table unless all demand and supply becomes zero

Thus, we obtain an initial basic feasible solution.

Tables below explain the Vogel Approximation Method with problem1

Table 2: Vogel Approximation Method

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M_1</th>
<th>M_2</th>
<th>M_3</th>
<th>M_4</th>
<th>Availability</th>
</tr>
</thead>
<tbody>
<tr>
<td>W_1</td>
<td>23</td>
<td>43</td>
<td>14(50)</td>
<td>10</td>
<td></td>
</tr>
<tr>
<td>W_2</td>
<td>32</td>
<td>12(80)</td>
<td>10</td>
<td></td>
<td></td>
</tr>
<tr>
<td>W_3</td>
<td>22</td>
<td>60</td>
<td>60</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Demand 70 Total=240

The total transportation cost is
23*10+14*50+12*80+20*40+22*60=4010

Penalty 4 1 9 26
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Step 2

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M₁</th>
<th>M₂</th>
<th>M₃</th>
<th>M₄</th>
<th>Availability</th>
<th>Penalty</th>
</tr>
</thead>
<tbody>
<tr>
<td>W₁</td>
<td>24</td>
<td>23</td>
<td>14(50)</td>
<td>10</td>
<td>9</td>
<td></td>
</tr>
<tr>
<td>W₂</td>
<td></td>
<td></td>
<td>12(80)</td>
<td></td>
<td>20 ←</td>
<td></td>
</tr>
<tr>
<td>W₃</td>
<td>20</td>
<td>22</td>
<td></td>
<td>100</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>Demand</td>
<td>40</td>
<td>70</td>
<td></td>
<td></td>
<td>9</td>
<td></td>
</tr>
</tbody>
</table>

Step 3

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M₁</th>
<th>M₂</th>
<th>M₃</th>
<th>M₄</th>
<th>Availability</th>
<th>Penalty</th>
</tr>
</thead>
<tbody>
<tr>
<td>W₁</td>
<td>23</td>
<td></td>
<td>14(50)</td>
<td>10</td>
<td>1</td>
<td></td>
</tr>
<tr>
<td>W₂</td>
<td></td>
<td>12(80)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>W₃</td>
<td>20(40)</td>
<td>22</td>
<td></td>
<td>60</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>Demand</td>
<td></td>
<td></td>
<td></td>
<td>70</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Penalty</td>
<td>4</td>
<td>1</td>
<td></td>
<td></td>
<td>9</td>
<td></td>
</tr>
</tbody>
</table>

Step 4

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M₁</th>
<th>M₂</th>
<th>M₃</th>
<th>M₄</th>
<th>Availability</th>
<th>Penalty</th>
</tr>
</thead>
<tbody>
<tr>
<td>W₁</td>
<td>23(10)</td>
<td></td>
<td>14(50)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>W₂</td>
<td></td>
<td></td>
<td>12(80)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>W₃</td>
<td>20(40)</td>
<td>22(60)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

The total transportation cost is
23*10+14*50+12*80+20*40+22*60=4010

From this cost it looks like that Vogel’s Approximation method as well as Least Cost Method is giving same basic feasible solution. But it is not true in most of the cases Vogel’s Approximation method gives better solution. Hence, we are considering one more example to show that VAM is better than Least Cost Method.

Problem 2

Solve the following transportation problem

<table>
<thead>
<tr>
<th>From</th>
<th>To</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>Supply</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2</td>
<td>7</td>
<td>4</td>
<td>5</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>3</td>
<td>3</td>
<td>1</td>
<td>8</td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>5</td>
<td>4</td>
<td>7</td>
<td>7</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>1</td>
<td>6</td>
<td>2</td>
<td>14</td>
<td></td>
</tr>
<tr>
<td>Demand</td>
<td>7</td>
<td>9</td>
<td>18</td>
<td>34</td>
<td></td>
</tr>
</tbody>
</table>

The total Transportation cost is
7*2+4*3+1*8+4*7+2*7+1*7=83

The table below gives the basic feasible solution obtained by using Vogel’s Approximation Method

<table>
<thead>
<tr>
<th>From</th>
<th>To</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>Supply</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2</td>
<td>7</td>
<td>4</td>
<td>5</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>3</td>
<td>3</td>
<td>1</td>
<td>8</td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>5</td>
<td>4</td>
<td>7</td>
<td>7</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>1</td>
<td>6</td>
<td>2</td>
<td>14</td>
<td></td>
</tr>
<tr>
<td>Demand</td>
<td>7</td>
<td>9</td>
<td>18</td>
<td>34</td>
<td></td>
</tr>
</tbody>
</table>

The total Transportation cost is
2*5+3*2+1*6+4*7+1*2+2*12= 76

Hence, from above problem, we can conclude the basic feasible solution obtained from Vogel’s Approximation is better Least Cost Method.

Modified Distribution Method

The basic feasible solution obtained by above methods are not necessarily optimal, therefore, we have to check the optimality of the above solution. The most reliable method to check the optimality of the basic feasible solution is Modified Distribution Method (MODI Method)

It’s a fact that solution obtained by VAM is more nearer to optimal solution than any other method. The necessary condition for applying this method is there must be m+n-1 allocations; otherwise it is a case of degeneracy, which we discuss later on.

Algorithm

Step 1: We find the initial basic feasible solution by applying any of the above methods, preferably with VAM on a balanced transportation problem.

Step 2: Assign variables $u_i$ and $v_j$ to each row and columns. Then $u_i$ or $v_j$ should be taken as zero, for which maximum number of allocations are made in its row or column.

Step 3: For allocated cells, we divide the cost as

$c_{ij} = u_i + v_j$

Step 4: Examine the sign for each $\Delta_{ij}$

i) If all $\Delta_{ij} \geq 0$, then the optimality criterion has been satisfied and the initial cost obtained is optimum. If some $\Delta_{ij}=0$, then alternate optimal solution exists.
ii) If some \( \Delta_{ij} \leq 0 \), then the cost obtained is not optimal. It has to be reduced by some technique (given in next step).

Step 5: In the cell having the most negative value of \( \Delta_{ij} \), we will make allocations. Rescheduling of allocations is done by the looping process.

Find the closed path with the selected unoccupied cell and allocate an unknown quantity \( \Theta \), to the cell. Add and subtract interchangeably, \( \Theta \) to and from the transition cells of the loop in such a way that the rim requirements remain satisfied.

Firstly we will take a solution obtained by least cost method and check its optimality, as we know the solution obtained by this method; it helps us to learn the whole algorithm.

Step1

\[
\begin{array}{cccc}
(2) & (3) & (1) & (4) & u_1=4 \\
(3) & (2) & (4) & (1) & u_2=1 \\
(5) & (0) & (5) & (7) & u_3=1 \\
(1) & 7 & (6) & (5) & u_4=2 \\
\end{array}
\]

Step 2

\[
\begin{array}{cccc}
(2) & (1) & (3) & (4) & u_1=4 \\
(3) & (1) & (3) & (2) & u_2=1 \\
(5) & (4) & (7) & (1) & u_3=1 \\
(1) & 7 & (6) & (5) & u_4=2 \\
\end{array}
\]

As two \( \Delta_{ij} \) are negative, hence solution obtained is not optimal. Thus, we modify the distribution of allocation to make the solution optimal.

Step 3

\[
\begin{array}{cccc}
(2) & (3) & (1) & (4) & u_1=4 \\
(3) & (2) & (4) & (1) & u_2=1 \\
(5) & (4) & (7) & (2) & u_3=1 \\
(1) & 7 & (6) & (5) & u_4=2 \\
\end{array}
\]

Now total cost = \( 3*2 + 4*5 + 4*7 + 1*6 + 1*7 + 2*7 = 81 \), which has reduced by 2 units

Step 4

\[
\begin{array}{cccc}
(2) & 5 & (1) & (4) & u_1=2 \\
(3) & 2 & 6 & 1 & u_2=0 \\
(5) & 7 & 2 & 5 & u_3=1 \\
(1) & 12 & 6 & 10 & u_4=1 \\
\end{array}
\]

Step 5

As all \( \Delta_{ij} \) are positive; hence solution obtained is optimal.

**Special Case of Degeneracy**

**Problem 3**

There are three factories and they are supplying their goods to 4 markets, the cost matrix of the problem is given below. Obtain the optimal transportation cost.

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M1</th>
<th>M2</th>
<th>M3</th>
<th>M4</th>
<th>Availability</th>
</tr>
</thead>
<tbody>
<tr>
<td>F1</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>4</td>
<td>6</td>
</tr>
<tr>
<td>F2</td>
<td>4</td>
<td>3</td>
<td>2</td>
<td>0</td>
<td>8</td>
</tr>
<tr>
<td>F3</td>
<td>0</td>
<td>2</td>
<td>2</td>
<td>1</td>
<td>10</td>
</tr>
</tbody>
</table>

| Demand | 4  | 6  | 8  | 6  | Total=24     |

Table below gives solution using Vogel’s Approximation Method for problem 3

<table>
<thead>
<tr>
<th>Factory / Markets</th>
<th>M1</th>
<th>M2</th>
<th>M3</th>
<th>M4</th>
<th>Availability</th>
</tr>
</thead>
<tbody>
<tr>
<td>F1</td>
<td>1</td>
<td>2(6)</td>
<td>3</td>
<td>4</td>
<td>6</td>
</tr>
<tr>
<td>F2</td>
<td>4</td>
<td>3</td>
<td>2(2)</td>
<td>0(6)</td>
<td>8</td>
</tr>
<tr>
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<td>0(4)</td>
<td>2</td>
<td>2(6)</td>
<td>1</td>
<td>10</td>
</tr>
</tbody>
</table>

| Demand | 4  | 6  | 8  | 6  | Total=24     |

Minimum cost = \( 2*6 + 2*2 + 0*6 + 0*4 + 2*6 = 28 \)

As we can see that number of allocations are 5 which are less than \( m+n-1 \) allocations i.e., \( 3+4-1=6 \), hence this is a problem of degeneracy, now to resolve this degeneracy and check its optimality using MODI (Modified Distribution Method), we will follow the following steps.
First we will follow the regular steps of MODI method, then the rows in which there are no allocations will be assigned a very small value $\varepsilon_1$, $\varepsilon_2$ etc, depending on the number of less allocations than $m+n-1$. Most of the practical aspects we will not be having more than two less allocations.

**Step 1**

<table>
<thead>
<tr>
<th></th>
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<td>72</td>
<td>94</td>
<td>19</td>
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</table>

In this first step, on applying MODI’s method we are not able to get the value of $u_1$, there for we will put $\varepsilon$ in first row or second column and obtain the values of $u_1$ and $v_2$.

**Step 2**

In this way, we can easily resolve the problem of degeneracy in transportation problem. In my view it will help students in resolution of degeneracy in much easier way. In case of two less allocations similarly, we can put $\varepsilon_1$, $\varepsilon_2$ in rows or columns for which we are not getting the value of the variables.

The solution obtain in the above problem is optimal solution as all the differences obtained positive or zero, as one difference is zero, hence alternate optimal solution is possible.

**Problem 4**

Consider the case of five factories and five markets with per unit transportation cost

<table>
<thead>
<tr>
<th>Factory/Market</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>Supply</th>
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<td>20</td>
<td>72</td>
<td>94</td>
<td>19</td>
<td>8</td>
</tr>
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</table>

Table below gives solution with Vogel’s Approximation Method

<table>
<thead>
<tr>
<th>Factory/Market</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
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<th>Supply</th>
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<td>20</td>
<td>72</td>
<td>94</td>
<td>19</td>
<td>19</td>
</tr>
</tbody>
</table>

Which gives only eight allocations; hence it is the case of degeneracy.

Tables below give the solution of problem 4 using MODI method.

<table>
<thead>
<tr>
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<td>(27)</td>
<td>(20)</td>
<td>(72)</td>
<td>(94)</td>
<td>(19)</td>
</tr>
</tbody>
</table>

$u_1 = 79$  $u_2 = 48$  $u_3 = 17$  $u_4 = 0$  $u_5 = 5$  

$V_1 = -1$  $V_2 = 15$  $V_3 = 24$  $V_4 = 19$  $V_5 = -8$
On applying MODI’s method both first and second row has not got any allocation. Putting $\varepsilon$ in second row having minimum cost in both the rows we obtain the following matrix.

### Step 2

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<td>(94) (9)</td>
<td>(71) (41) (46) (4) (85)</td>
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<td>(57)</td>
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<td>(-1) (33)</td>
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<td>(72)</td>
<td>(94) (24)</td>
<td>(19)</td>
<td>(23) (29) (72) (70) (22)</td>
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</tbody>
</table>

$\mathbf{v}_1 = -1$  
$\mathbf{v}_2 = 15$  
$\mathbf{v}_3 = 24$  
$\mathbf{v}_4 = 19$  
$\mathbf{v}_5 = -8$

### Step 3

Shifting $\varepsilon$ to the cell having most negative element in first row, and checking the optimality using Modified Distribution Method, we have

<table>
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<tr>
<th></th>
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<td>(94) (24)</td>
<td>(19)</td>
<td>(23) (29) (72) (70) (22)</td>
</tr>
</tbody>
</table>

$\mathbf{v}_1 = 44$  
$\mathbf{v}_2 = 60$  
$\mathbf{v}_3 = 103$  
$\mathbf{v}_4 = 64$  
$\mathbf{v}_5 = 71$

### Step 4

Again modifying the distribution we obtain the reduced matrix with 9 allocations, hence degeneracy is resolved

<table>
<thead>
<tr>
<th></th>
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<th>2</th>
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<th>$\boldsymbol{U}_1$ = 0</th>
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$\mathbf{v}_1 = 44$  
$\mathbf{v}_2 = 72$  
$\mathbf{v}_3 = 103$  
$\mathbf{v}_4 = 64$  
$\mathbf{v}_5 = 71$
Step 5
Again modify the distribution we obtained the following matrix

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<td>(20) (6)</td>
<td>(72) (61)</td>
<td>(94) (22)</td>
<td>(19) (2)</td>
</tr>
</tbody>
</table>

\[ v_1 = 44 \quad v_2 = 62 \quad v_3 = 103 \quad v_4 = 64 \quad v_5 = 61 \]

The solution obtain in the above problem is optimal solution as all the differences obtained positive, hence, the total transportation cost comes out to be $2652

Special Cases

a) Unbalanced Transportation problem. A necessary and sufficient condition for existence of feasible solution is that total demand must be equal to the total supply, i.e.

\[ \sum_{i=1}^{m} a_i = \sum_{j=1}^{n} b_j \]

However, sometimes there may be more demand than the supply and vice versa in which case it is said to be unbalanced. It may occur in the following situations:

i) \[ \sum_{i=1}^{m} a_i > \sum_{j=1}^{n} b_j \] and ii) \[ \sum_{i=1}^{m} a_i < \sum_{j=1}^{n} b_j \]

In case I, we introduce a dummy destination in the transportation table. The cost of transporting to this are all set equal to zero. The requirement at this dummy destination is then assumed to be equal to

\[ \sum_{i=1}^{m} a_i - \sum_{j=1}^{n} b_j \]

In case II, we introduce a dummy source in the transportation table. The cost of transporting from this source to any destination is all set equal to zero. The availability at this dummy source is assumed to be equal to

\[ \sum_{j=1}^{n} b_j - \sum_{i=1}^{m} a_i \]

b) Maximization Problem. In general, transportation method is used for minimization problems. However, it can also be used to solve problem in which objective is to maximize when we consider the unit profit (or payoff) \( p_{ij} \) instead of unit cost \( c_{ij} \) associated with route \((i,j)\). The solution process to solve such problems is given below.

First, convert the given problem into minimization problem by replacing each element of the transportation table by its difference from the maximum element of the table.

Second, obtain an initial feasible solution using any of the three methods.

Third, use Modified Distribution Method for finding an optimal solution.

c) Prohibited Routes- The situation may arise such as road hazards (snow, flood, etc.), traffic regulations, strike etc., when it is not possible to transport goods from certain sources to certain destinations. Such type of problem can be handled by assigning infinite cost \((\infty)\) to that route (or cell)

d) Alternative optimal solutions- the existence of alternative optimal solutions can be determined by an inspection of the opportunity costs, \( z_{ij} - c_{ij} \) for the unoccupied cells. If an unoccupied cells. If an unoccupied cell in an optimal solution has opportunity cost of zero, then an alternative optimal solution can be formed with another set of allocations without increasing the total transportation cost. To obtain an alternate solution, trace a closed loop beginning with cell having \( \Delta_{ij} = 0 \), and get the revised solution in the same way as a solution is improved. It may be observed that this revised solution would also entail the same total cost as before. It goes without saying that for every ‘zero’ value of \( \Delta_{ij} \) value in the optimal solution tableau, a revised solution is obtained.

e) Assignment Problem- Assignment problem is a special case of transportation problem having equal number of sources and destinations (sink), and all \( a_i \) and \( b_j \) equal to one. Thus, assignment problem can also solve with the help transportation algorithm.

f) Transshipment Problem- A transportation problem can be converted into a transshipment problem by relaxing the restrictions on the receiving and sending
units on the origins and destinations respectively. An 
m-origin, n-destination, transportation problem, 
when expressed as a transshipment problem, shall 
become an enlarged problem: with \( m + n \) origins and 
an equal number of destinations. With minor 
modification this problem can be solved using the 
transportation algorithm.

2. Conclusions

Transportation problem is one of most popular type of 
linear programming problem. The algorithm to get solution 
of transportation problem has been widely used to solve 
transportation problem in various cases. The only drawback 
with transportation problem is the case of degeneracy i.e. 
when numbers of allocations are less than \( m+n-1 \), there was 
no exact rule to resolve this degeneracy. In this study, a 
method is developed to resolve the problem of degeneracy. 
Hence, this study tries to make transportation algorithm 
more efficient and user friendly.

Appendix

Theorem 1 (Existence of a feasible solution)

The necessary and sufficient condition for the existence of 
a feasible solution to the Transportation Problem is 
\[
\sum_{i=1}^{m} a_i = \sum_{j=1}^{n} b_j ,
\]
that is, the total capacity (supply) must 
equal total requirement (demand)

Theorem 2

The dimensions of the basis of a Transportation Problem 
are \((m+n-1) \times (m+n-1)\). This means that a Transportation 
problem has only \((m+n-1)\)-independent structural constraints 
and its basic feasible solution has only \((m+n-1)\)-positive 
components.

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